

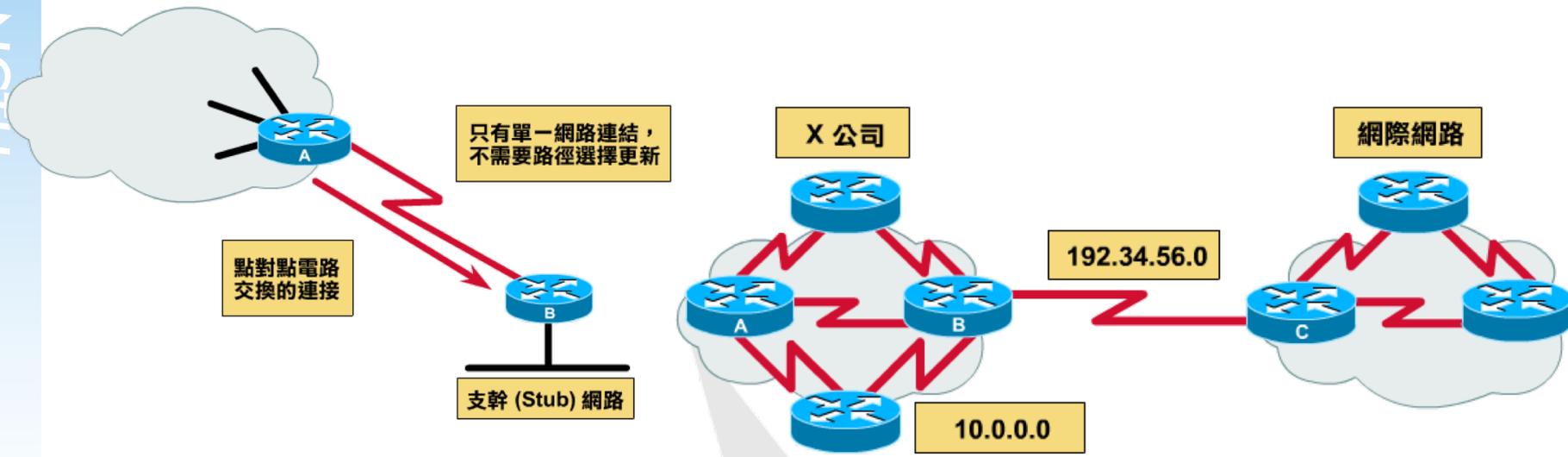
Routing

jnlin

Why dynamic route ? (1)

❑ Static route is ok only when

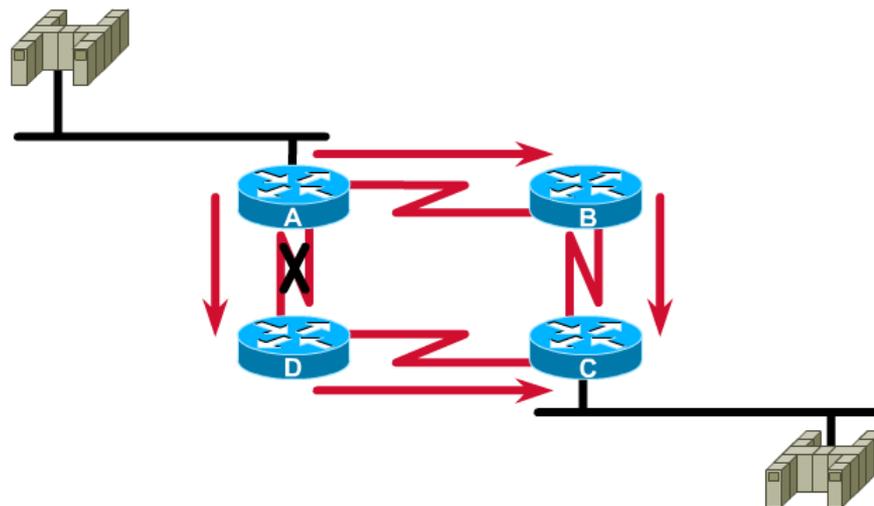
- Network is small
- There is a single connection point to other network
- No redundant route



Why dynamic route ? (2)

□ Dynamic Routing

- Routers update their routing table with the information of adjacent routers
- Dynamic routing need a routing protocol for such communication
- Advantage:
 - They can react and adapt to changing network condition



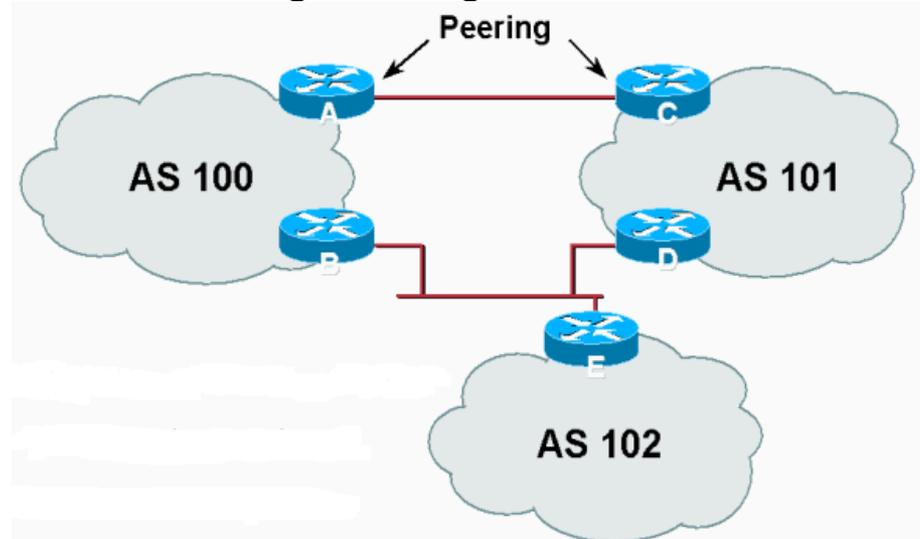
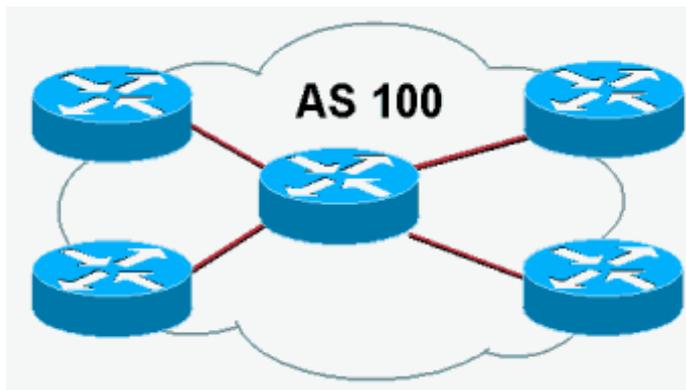
Routing Protocol

- ❑ Used to change the routing table according to various routing information
 - Specify detail of communication between routers
 - Specify information changed in each communication,
 - Network reachability
 - Network state
 - Metric
- ❑ Metric
 - A measure of how good a particular route
 - Hop count, bandwidth, delay, load, reliability, ...
- ❑ Each routing protocol may use different metric and exchange different information

Autonomous System

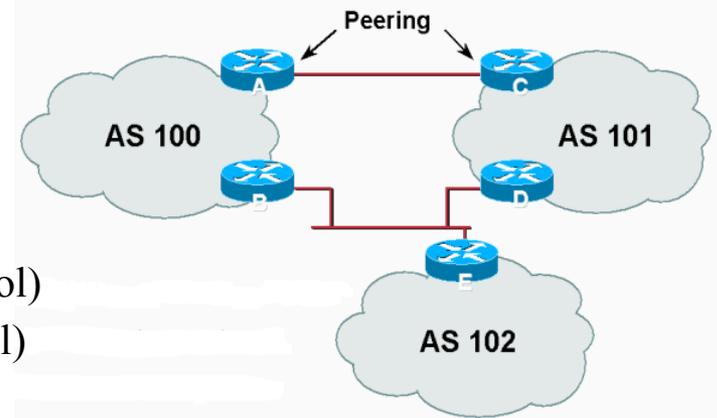
❑ Autonomous System (AS)

- Internet is organized into a collection of autonomous system
- An AS is a collection of networks with same routing policy
 - Single routing protocol
 - Normally administered by a single entity
 - Corporation or university campus
 - All depend on how you want to manage routing

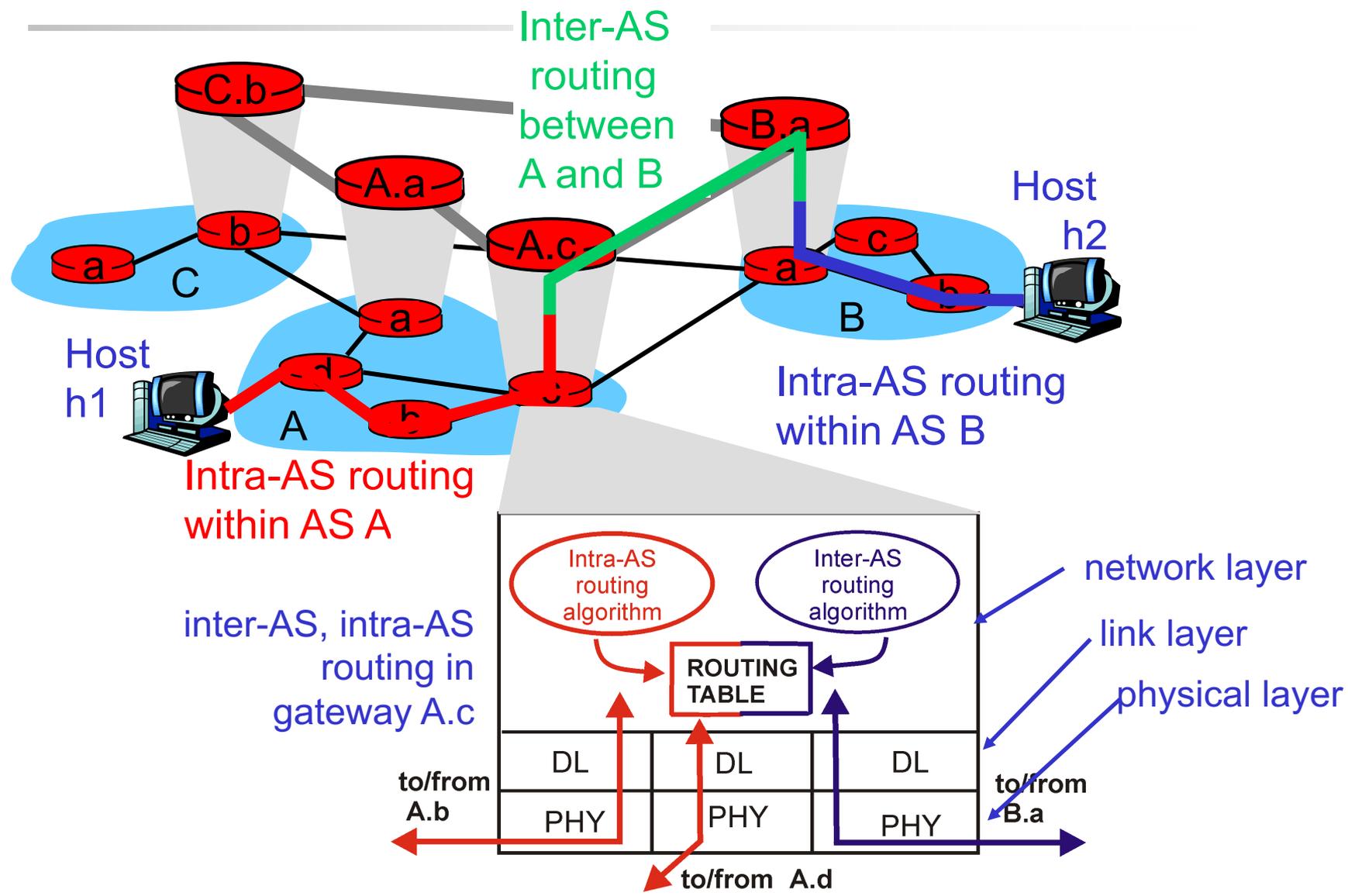


Category of Routing Protocols – by AS

- ❑ AS-AS communication
 - Communications between routers in different AS
 - Interdomain routing protocols
 - Exterior gateway protocols (EGP)
 - Ex:
 - BGP (Border Gateway Protocol)
- ❑ Inside AS communication
 - Communication between routers in the same AS
 - Intradomain routing protocols
 - Interior gateway protocols (IGP)
 - Ex:
 - RIP (Routing Information Protocol)
 - IGRP (Interior Gateway Routing Protocol)
 - OSPF (Open Shortest Path First Protocol)



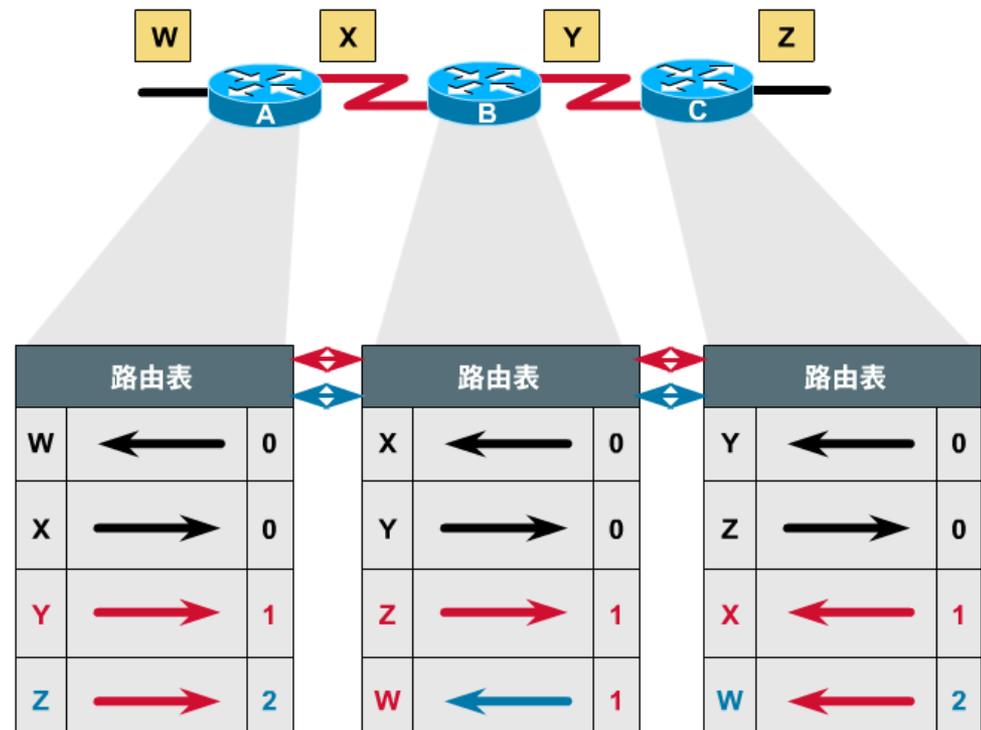
Intra-AS and Inter-AS routing



Category of Routing Protocols – by information changed (1)

□ Distance-Vector Protocol

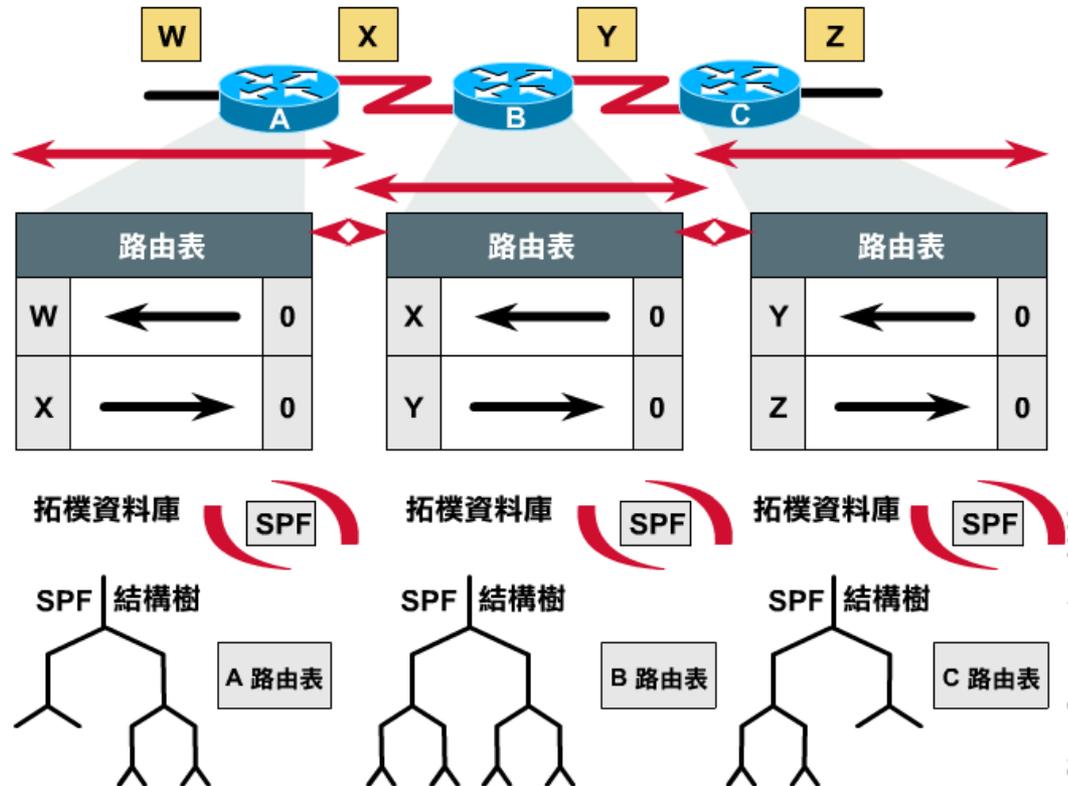
- Message contains a vector of distances, which is the cost to other network
- Each router updates its routing table based on these messages received from neighbors
- Protocols:
 - RIP
 - IGRP
 - BGP



Category of Routing Protocols – by information changed (2)

□ Link-State Protocol

- Broadcast their link state to neighbors and build a complete network map at each router using Dijkstra algorithm
- Protocols:
 - OSPF

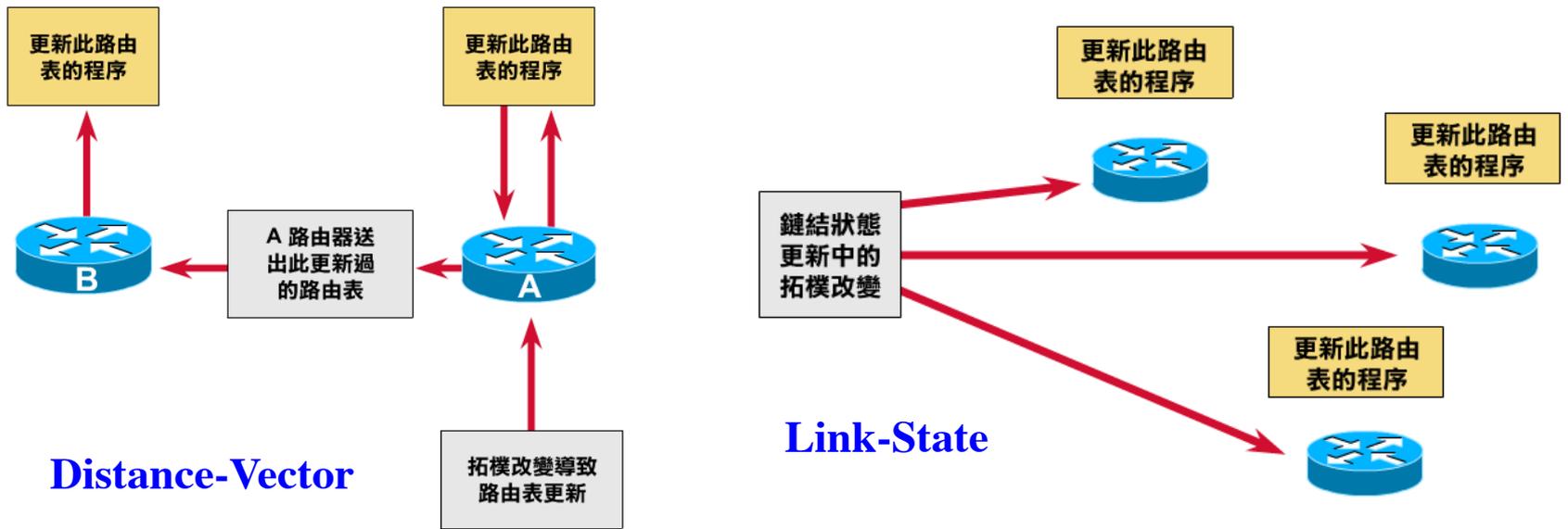


Difference between Distance-Vector and Link-State

□ Difference

	Distance-Vector	Link-State
Update	updates neighbor (propagate new info.)	update all nodes
Convergence	Propagation delay cause slow convergence	Fast convergence
Complexity	simple	Complex

□ Information update sequence



Routing Protocols

RIP	IGP, DV
IGRP	IGP, DV
OSPF	IGP, LS
BGP	EGP

RIP

❑ RIP

- Routing Information Protocol

❑ Category

- Interior routing protocol
- Distance-vector routing protocol
 - Using "hop-count" as the cost metric

❑ Example of how RIP advertisements work

Destination network	Next router	# of hops to destination
1	A	2
20	B	2
30	B	7

Routing table in router before
Receiving advertisement

Destination network	Next router	# of hops to destination
30	C	4
1	--	1
10	--	1

Advertisement from router A

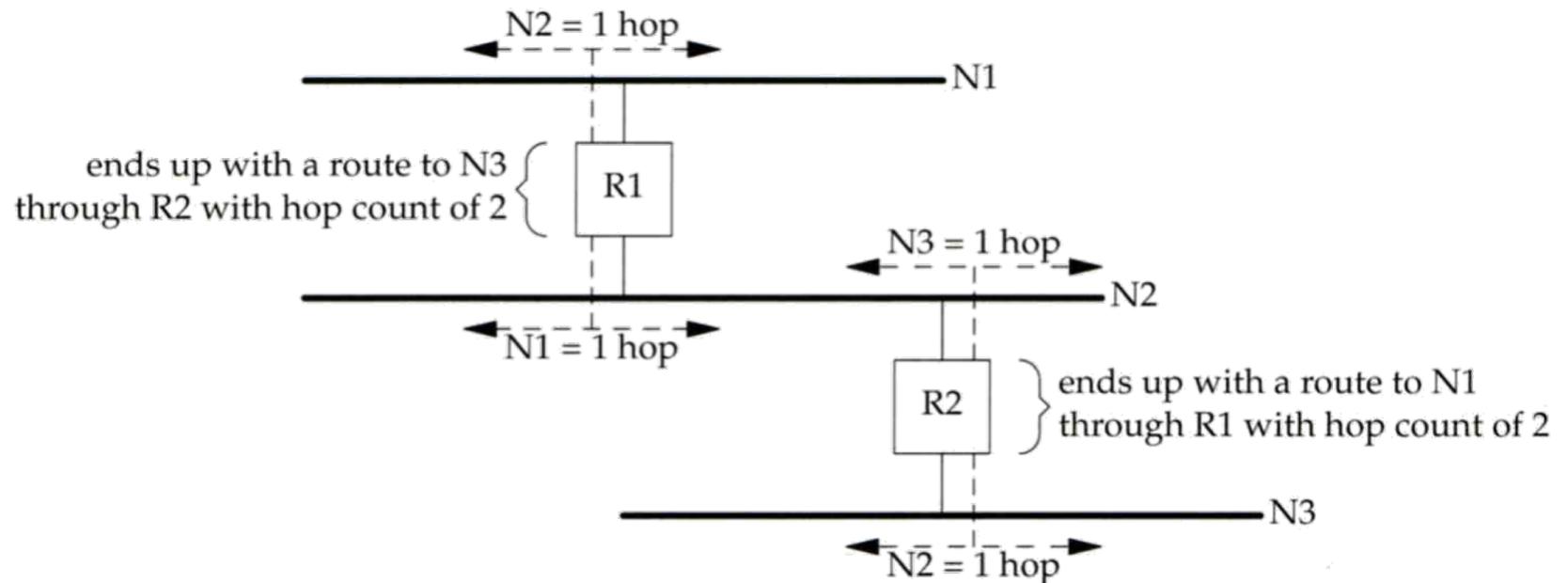
Destination network	Next router	# of hops to destination
1	A	2
20	B	2
30	A	5

Routing table after
receiving advertisement

RIP

– Example

□ Another example

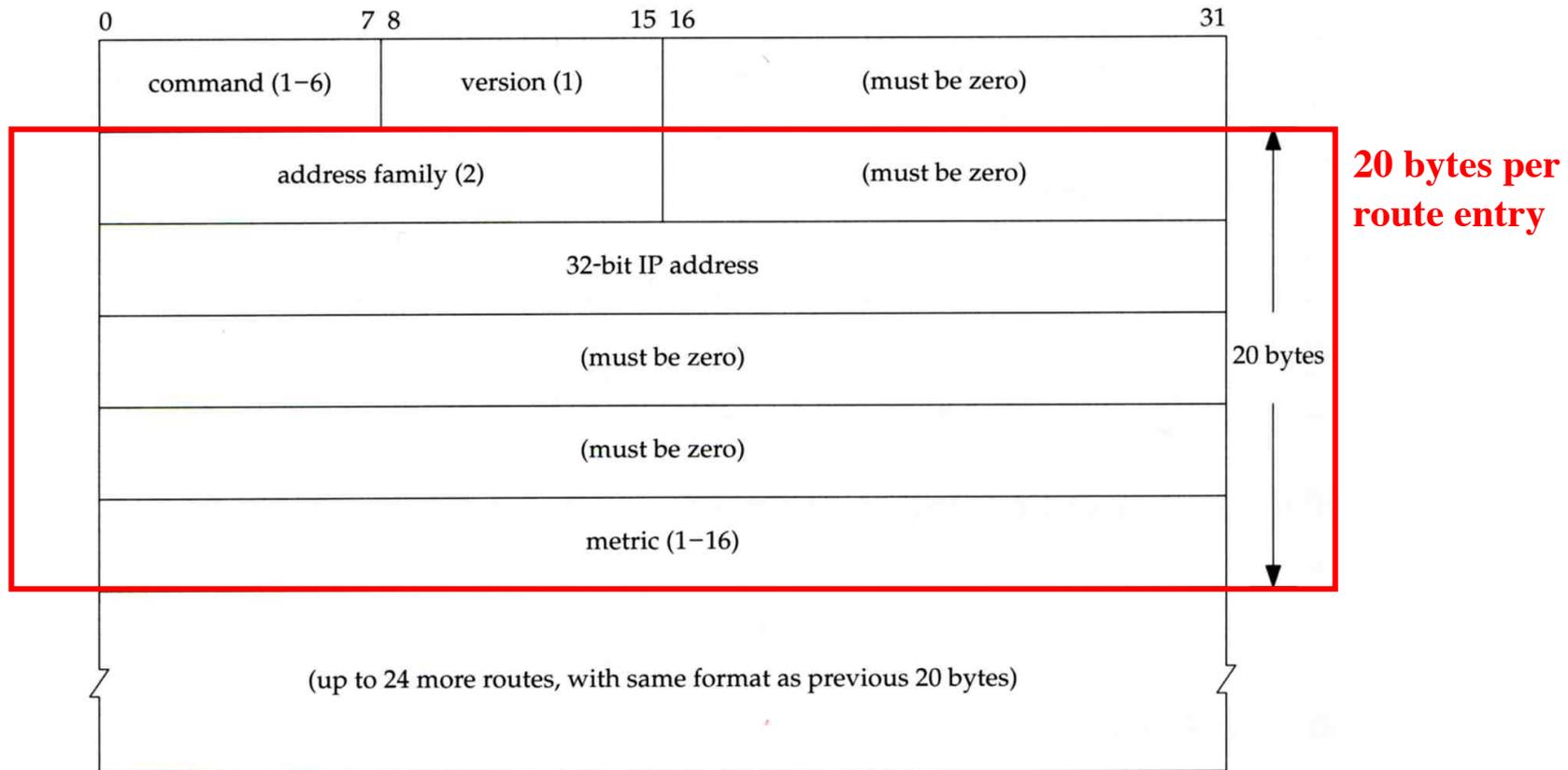


RIP

– Message Format

❑ RIP message is carried in UDP datagram

- Command: 1 for request and 2 for reply
- Version: 1 or 2 (RIP-2)



RIP

– Operation

- ❑ routed – RIP routing daemon
 - Operated in UDP port 520
- ❑ Operation
 - Initialization
 - Probe each interface
 - send a request packet out each interface, asking for other router's complete routing table
 - Request received
 - Send the entire routing table to the requestor
 - Response received
 - Add, modify, delete to update routing table
 - Regular routing updates
 - Router sends out their routing table to every neighbor every 30 seconds
 - Triggered updates
 - Whenever a route entry's metric change, send out those changed part routing table

RIP

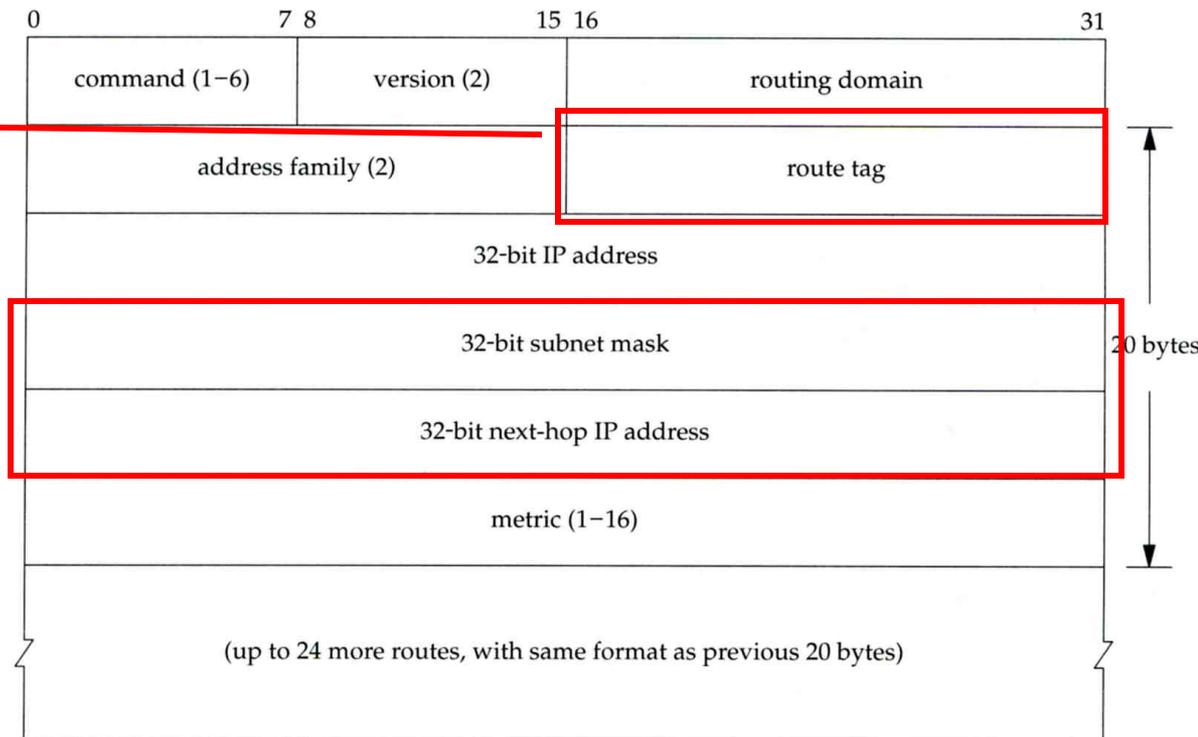
– Problems of RIP

❑ Issues

- 15 hop-count limits
- Take long time to stabilize after the failure of a router or link
- No CIDR

❑ RIP-2

- EGP support
 - AS number
- CIDR support



IGRP (1)

- ❑ IGRP – Interior Gateway Routing Protocol
- ❑ Similar to RIP
 - Interior routing protocol
 - Distance-vector routing protocol
- ❑ Difference between RIP
 - Complex cost metric other than hop count
 - delay time, bandwidth, load, reliability
 - The formula

$$\left(\frac{\textit{bandwidth_weight}}{\textit{bandwidth} * (1 - \textit{load})} + (\textit{delay_weight} * \textit{delay}) \right) * \textit{reliability}$$

- Use TCP to communicate routing information
- Cisco System's proprietary routing protocol

IGRP (2)

- ❑ Advantage over RIP
 - Control over metrics
- ❑ Disadvantage
 - Still classful and has propagation delay

OSPF (1)

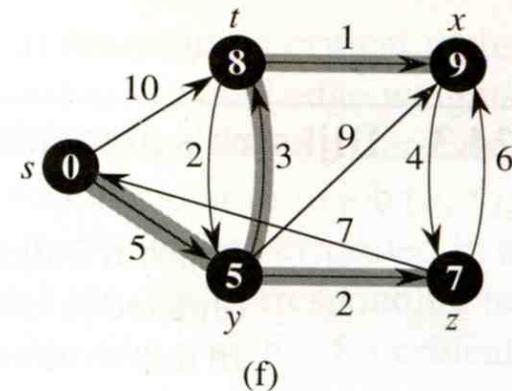
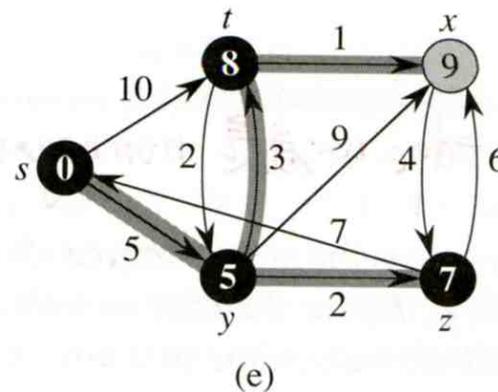
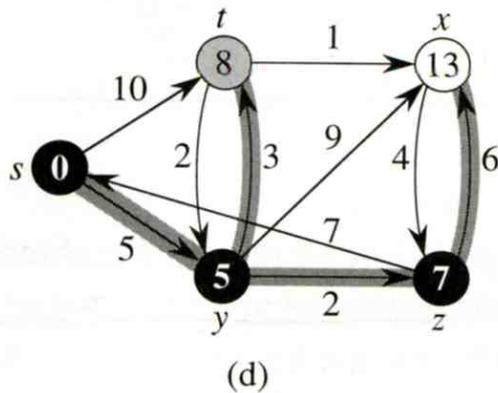
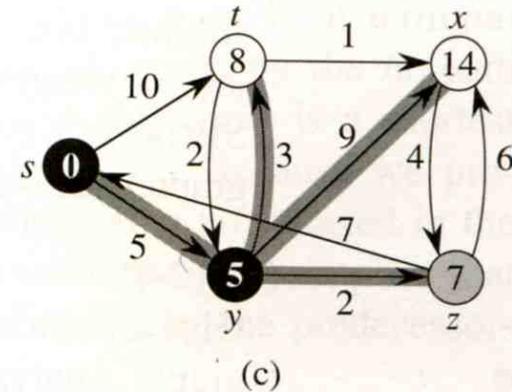
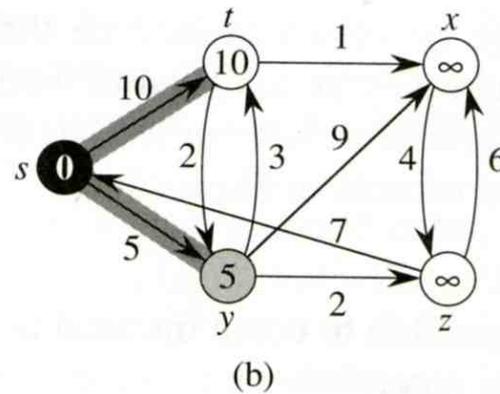
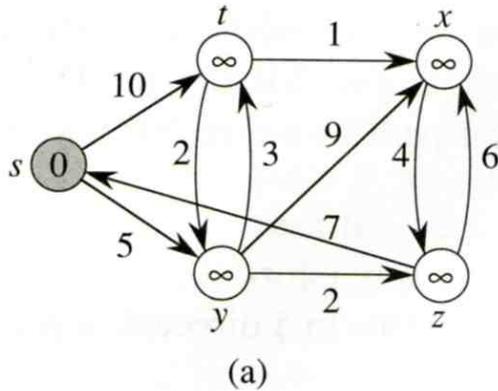
- ❑ OSPF
 - Open Shortest Path First
- ❑ Category
 - Interior routing protocol
 - Link-State protocol
- ❑ Each interface is associated with a cost
 - Generally assigned manually
 - The sum of all costs along a path is the metric for that path
- ❑ Neighbor information is broadcast to all routers
 - Each router will construct a map of network topology
 - Each router run Dijkstra algorithm to construct the shortest path tree to each routers

OSPF

– Dijkstra Algorithm

□ Single Source Shortest Path Problem

- Dijkstra algorithm use “greedy” strategy
- Ex:

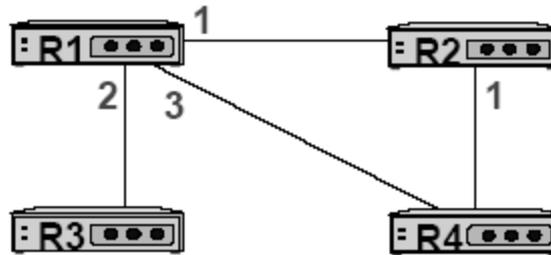


OSPF

- Routing table update example (1)

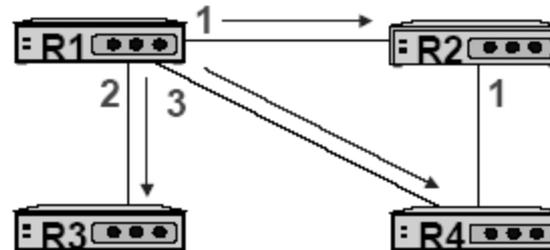
R1

D	Path	M
R1		
R2		
R3		
R4		



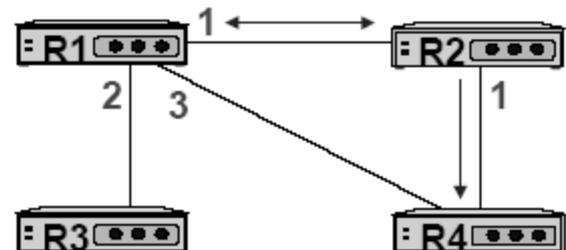
R1

D	Path	M
R1	direct	0
R2	R1-R2	1
R3	R1-R3	2
R4	R1-R4	3



R1

D	Path	M
R1	direct	0
R2	R1-R2	1
R3	R1-R3	2
R4	R1-R4	3

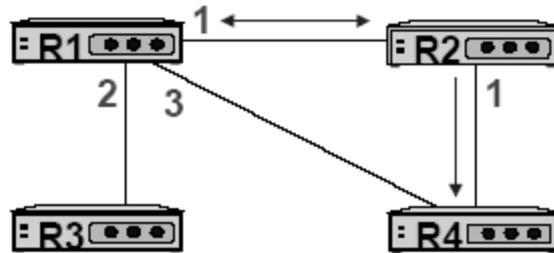


OSPF

- Routing table update example (2)

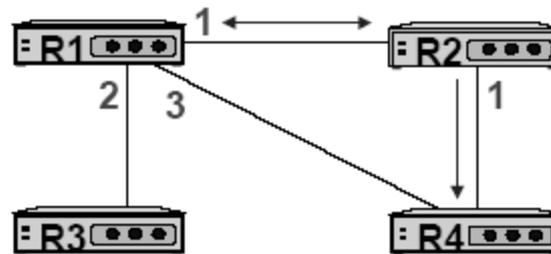
R1

D	Path	M
R1	direct	0
R2	R1-R2	1
R3	R1-R3	2
R4	R1-R4	3



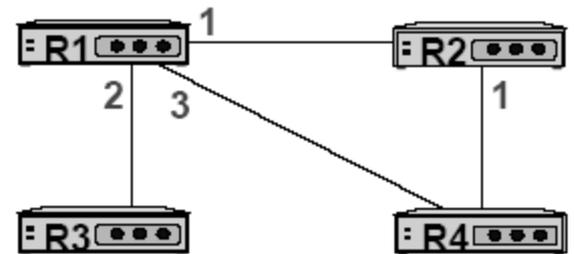
R1

D	Path	M
R1	direct	0
R2	R1-R2	1
R3	R1-R3	2
R4	R1-R2-R4	2



R1

D	Path	M
R1	direct	0
R2	R1-R2	1
R3	R1-R3	2
R4	R1-R2-R4	2



OSPF

– Summary

❑ Advantage

- Fast convergence
- CIDR support
- Multiple routing table entries for single destination, each for one type-of-service
 - Load balancing when cost are equal among several routes

❑ Disadvantage

- Large computation

BGP

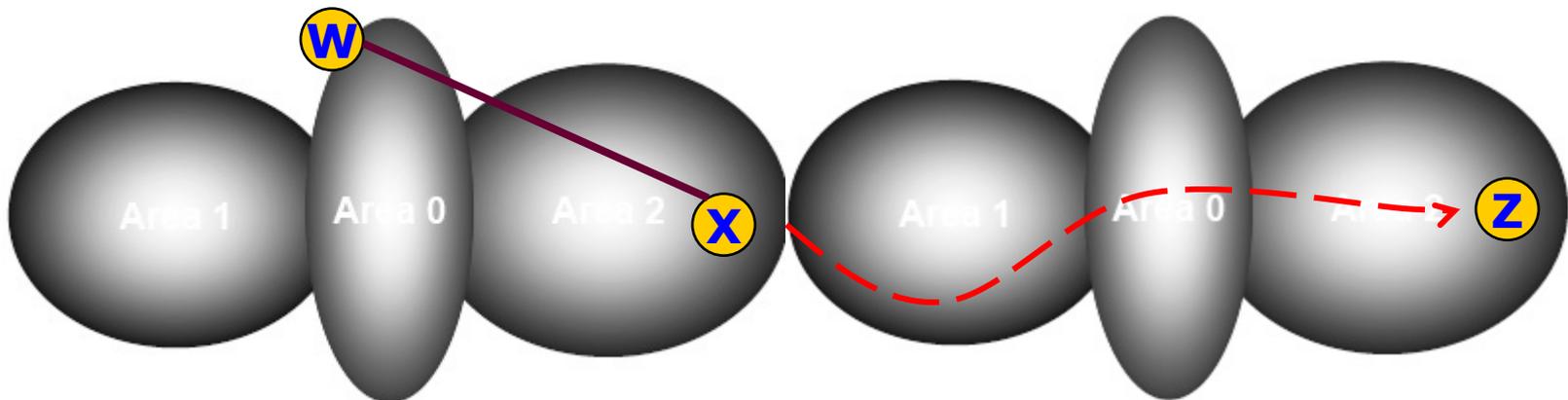
- ❑ BGP
 - Border Gateway Protocol
- ❑ Exterior routing protocol
 - Now BGP-4
 - Exchange network reachability information with other BGP systems
- ❑ Routing information exchange
 - Message:
 - Full path of autonomous systems that traffic must transit to reach destination
 - Can maintain multiple route for a single destination
 - Exchange method
 - Using TCP
 - Initial: entire routing table
 - Subsequent update: only sent when necessary
 - Advertise only optimal path
- ❑ Route selection
 - Shortest AS path

BGP

– Operation Example

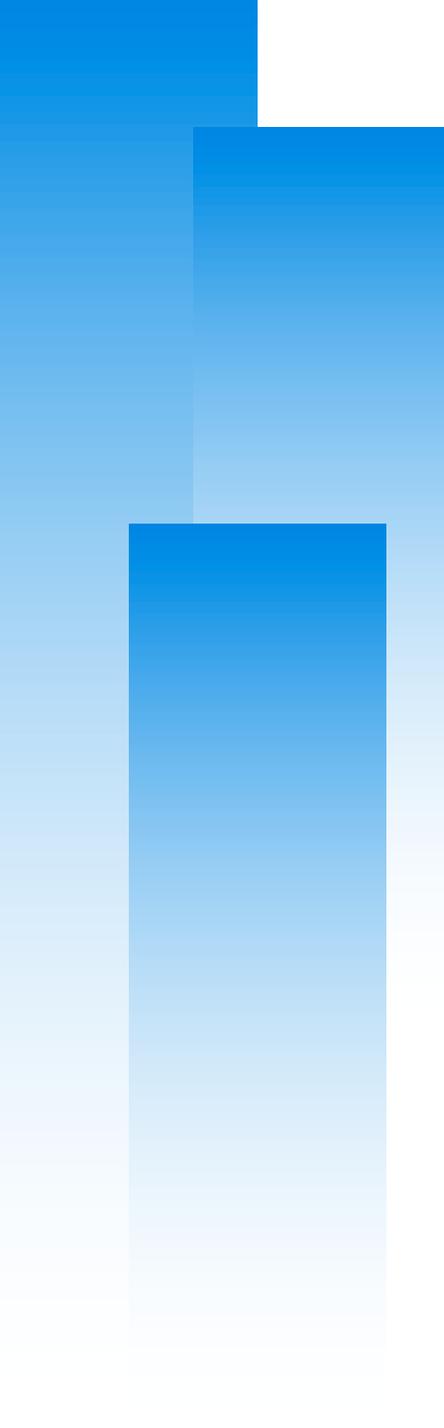
□ How BGP work

- The whole Internet is a graph of autonomous systems
- $X \rightarrow Z$
 - Original: $X \rightarrow A \rightarrow B \rightarrow C \rightarrow Z$
 - X advertise this best path to his neighbor W
- $W \rightarrow Z$
 - $W \rightarrow X \rightarrow A \rightarrow B \rightarrow C \rightarrow Z$



Routing Protocols Comparison

	RIP	IGRP	OSPF	BGP4
DV or LS	DV	DV	LS	Path Vec
TCP/UDP & Port	U - 520	IP - 9	T - 89	T - 179
Classless	No	No	Yes	Yes
Updates	Per.	Per.	Both	Trig.
Load Balance	No	Yes	Yes	No
Internal / External	Int.	Int.	Int.	Ext.
Metric	Hop Count	Load Errors Delay Bdwth	Sum of Int. Cost	Short. AS Path



routed

routed

❑ Routing daemon

- Speak RIP (v1 and v2)
- Supplied with most every version of UNIX
- Two modes
 - Server mode (-s) & Quiet mode (-q)
 - Both listen for broadcast, but server will distribute their information
- routed will add its discovered routes to kernel's routing table
- Support configuration file - /etc/gateways
 - Provide static information for initial routing table

```
net Nname[/mask] gateway Gname metric value <passive | active | extern>
host Hname gateway Gname metric value <passive | active | extern>
```